

Error Detection**Q1.**

- a. In a CRC error detection scheme, choose  $P(x) = X^4 + X + 1$ . Encode the bits 10010011011 using Shifted Polynomial method.

**Converting data to polynomial format;**

$$10010011011 = X^{10} + X^7 + X^4 + X^3 + X + 1$$

**Establishing length of FCS;**

$$P(x) = X^4 + X + 1 = 10011 \rightarrow \text{So length of FSC will be } 5 \text{ bits} - 1 = 4 \text{ bits.}$$

**Applying shifting to data;**

$$X^4(X^{10} + X^7 + X^4 + X^3 + X + 1) = X^{14} + X^{11} + X^8 + X^7 + X^5 + X^4$$

**Calculate FCS;**

$$\begin{array}{r}
 X^{10} + X^6 + X^4 + X^2 \\
 \hline
 X^4 + X + 1 \quad / \quad X^{14} + X^{11} + X^8 + X^7 + X^5 + X^4 \\
 \quad \quad \quad X^{14} + X^{11} + X^{10} \\
 \quad \quad \quad \hline
 \quad \quad \quad X^{10} + X^8 + X^7 + X^5 + X^4 \\
 \quad \quad \quad X^{10} + \quad \quad X^7 + X^6 \\
 \quad \quad \quad \hline
 \quad \quad \quad \quad X^8 + X^6 + X^5 + X^4 \\
 \quad \quad \quad \quad X^8 + \quad \quad X^5 + X^4 \\
 \quad \quad \quad \quad \hline
 \quad \quad \quad \quad \quad X^6 \\
 \quad \quad \quad \quad \quad X^6 + X^3 + X^2 \\
 \quad \quad \quad \quad \quad \hline
 \quad \quad \quad \quad \quad \quad X^3 + X^2
 \end{array}$$

**So;**

$$FSC = X^3 + X^2$$

**Converting to bits;**

$$FSC = X^3 + X^2 = 1100$$

**Append to data bits;**

$$\text{Bits sent} = 100100110111100$$

**Note :**

**Highlight to students when to stop this process. It stops when the power of the leading X is SMALLER than the leading power of the pattern.**

- b. Repeat the above problem using Base-2 method.

**Using Base 2 method****In base 2 method, we use XOR between the data bits and the pattern bits.**

**XOR table**

A	B	A XOR B
0	0	0
0	1	1
1	0	1
1	1	0

**Data bits are;**  
**10010011011**

**Establishing length of FCS;**

$P(x) = X^4 + X + 1 = 10011$  -> So length of FSC will be 5 bits – 1 = 4 bits.

**Shifting to left and inserting 4 zeros in the end (similar to shifted polynomial)**

**100100110110000** ←

**Calculate FCS;**

**100100110110000**

**10011**

-----

**10110110000**

**10011**

-----

**101110000**

**10011**

-----

**1000000**

**10011**

-----

**1100**

**So;**

**FCS = 1100**

**Append to data bits;**

**Bits sent = 100100110111100**

- c. Suppose the channel introduces errors and the string 000110110111100 is received. Can the error be detected?

**Bits received = 000110110111100**

$= X^{11} + X^{10} + X^8 + X^7 + X^5 + X^4 + X^3 + X^2$  *Note to students ; No shifting done on receive side !*

$$\begin{array}{r}
 \mathbf{X}^7 + \mathbf{X}^6 + \mathbf{X}^3 + \mathbf{X}^2 + \mathbf{X} \\
 \hline
 \mathbf{X}^4 + \mathbf{X} + \mathbf{1} \quad / \quad \begin{array}{r} \mathbf{X}^{11} + \mathbf{X}^{10} + \mathbf{X}^8 + \mathbf{X}^7 + \mathbf{X}^5 + \mathbf{X}^4 + \mathbf{X}^3 + \mathbf{X}^2 \\ \mathbf{X}^{11} + \mathbf{X}^8 + \mathbf{X}^7 \end{array} \\
 \hline
 \begin{array}{r} \mathbf{X}^{10} + \mathbf{X}^5 + \mathbf{X}^4 + \mathbf{X}^3 + \mathbf{X}^2 \\ \mathbf{X}^{10} + \mathbf{X}^7 + \mathbf{X}^6 \end{array} \\
 \hline
 \begin{array}{r} \mathbf{X}^7 + \mathbf{X}^6 + \mathbf{X}^5 + \mathbf{X}^4 + \mathbf{X}^3 + \mathbf{X}^2 \\ \mathbf{X}^7 + \mathbf{X}^4 + \mathbf{X}^3 \end{array} \\
 \hline
 \begin{array}{r} \mathbf{X}^6 + \mathbf{X}^5 + \mathbf{X}^2 \\ \mathbf{X}^6 + \mathbf{X}^3 + \mathbf{X}^2 \end{array} \\
 \hline
 \begin{array}{r} \mathbf{X}^5 + \mathbf{X}^3 \\ \mathbf{X}^5 + \mathbf{X}^2 + \mathbf{X} \end{array} \\
 \hline
 \mathbf{X}^3 + \mathbf{X}^2 + \mathbf{X}
 \end{array}$$

**ERRORS DETECTED...!!!!**

### Base-2 Method:

**Received bits**  
**000110110111100**

### Perform base 2 method again

```

000110110111100
10011
-----
100000110111100
10011
-----
      110110111100
      10011
      -----
        10000111100
        10011
        -----
          11111100
          10011
          -----
            1100100
            10011
            -----
              101000

```

$$\begin{array}{r} 10011 \\ - 1110 \\ \hline \end{array}$$

**Since the remainder is not 0, there are errors detected.**

d. If 000010110111100 is received will an error be detected?

**Bits received** = 000010110111100 =  $X^{10} + X^8 + X^7 + X^5 + X^4 + X^3 + X^2$

**Do error check;**

$$\begin{array}{r}
 \text{no error check,} \\
 \mathbf{X^6 + X^4 + X^2} \\
 \hline
 \mathbf{X^4 + X + 1} \quad / \quad \mathbf{X^{10} + X^8 + X^7 + X^5 + X^4 + X^3 + X^2} \\
 \mathbf{X^{10} + \phantom{X^8} X^7 + X^6} \\
 \hline
 \mathbf{X^8 + X^6 + X^5 + X^4 + X^3 + X^2} \\
 \mathbf{X^8 \phantom{+ X^6} + X^5 + X^4} \\
 \hline
 \mathbf{X^6 + X^3 + X^2} \\
 \mathbf{X^6 + X^3 + X^2} \\
 \hline
 \mathbf{\phantom{X^6 + X^3 + X^2}} \\
 \mathbf{\phantom{X^6 + X^3 + X^2}}
 \end{array}$$

**NO ERRORS DETECTED (So CRC is not fool-proof)**

**Q2.**

What is the difference between odd parity and even parity?

**In parity error control, an extra bit is added to each row of data.**

**The value (1 or 0) of this extra bit is set according to whether odd or even parity is used.**

**In even parity, the bit is set so that the number of 1s in the row is even. EG**

Data bits						Parity bit
Data Row 1	1	1	0	1	0	1
Data Row 2	0	1	1	1	1	0

**Note that for even parity, the parity bit can be 1 or 0, as required.**

**In odd parity, the bit is set so that the number of 1s in the row is odd. EG**

Data bits						Parity bit
Data Row 1	1	1	0	1	0	0
Data Row 2	0	1	1	1	1	1

**Note that for odd parity, the parity bit can be 1 or 0, as required.**

**Two dimensional parity checking is also possible by adding parity bits to the columns as well.**